## Trial tiling implementation using tra\_ldf\_iso

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#### Overview

A trial approach to 2D tiling (over *i* and *j*) is implemented for the *tra\_ldf\_iso* subroutine.

This document describes an implementation using public variables and an alternative approach using derived types. A number of potential issues are then described.

## **Test configuration**

A 1° configuration of GYRE is used with 1 CPU. XIOS is not used.

#### **Branches**

- Public variables
- Derived types

### **Variables**

- Global variables
  - o ntsi, ntsj- start index of tile
  - o ntei, ntej- end index of tile
  - o ntsim1, ntsjm1- start index of tile, minus 1
  - o nteip1, ntejp1- end index of tile, plus 1
  - o ntile- tile number
- Parameters
  - o jpnitile, jpnjtile, jpnijtile- number of tiles
- Loop indices
  - o jtile- loop over tiles
- Namelist
  - o In\_tile- Logical control on use of tiling
  - o nn\_tile\_i, nn\_tile\_j- tile length
- Pre-processor macros
  - o IND\_2D- substitution for ALLOCATE or DIMENSION arguments
- Working variables
  - o iitile, ijtile- tile number
- · Dummy arguments
  - o kntile (ntile)

## Namelist and tiling decomposition

The tiling decomposition is defined by parameters *jpnitile*, *jpnjtile* and *jpnijtile*, which are analogous to the MPP decomposition parameters *jpni*, *jpnj* and *jpnij*.

These parameters are declared in par\_oce and set using a new namelist section, namtile.

!					 	
&namtile	!	parameters	of the	tiling		
!					 	

```
ln_tile = .false. ! Use tiling (T) or not (F)
nn_tile_i = 10 ! Length of tiles in i
nn_tile_j = 10 ! Length of tiles in j
/
```

The decomposition parameters are then calculated in dom\_nam.

```
IF( ln_tile ) THEN
    jpnitile = (jpi - 2) / nn_tile_i
    jpnjtile = (jpj - 2) / nn_tile_j
    IF( MOD( jpi - 2, nn_tile_i ) /= 0 ) jpnitile = jpnitile + 1
    IF( MOD( jpj - 2, nn_tile_j ) /= 0 ) jpnjtile = jpnjtile + 1
ELSE
    jpnitile = 1
    jpnjtile = 1
ENDIF
jpnijtile = jpnitile * jpnjtile
```

# Implementation of tile subdomains

The DO loop macros in *do\_loop\_substitute.h90* are modified to refer to a set of domain indices, instead of those for the full domain. A macro is added for *ALLOCATE* and *DIMENSION* arguments.

```
#define __kIs_
#define __kJs_
                  2
#define kIsm1
                 1
#define kJsm1
                  1
#define __kIs_
                 ntsi
#define __kJs_
                 ntsj
#define __kIsm1_
                 ntsim1
#define __kJsm1_
                 ntsjm1
#define __kIe_
                 jpim1
#define __kJe_
                 jpjm1
#define kIep1
                 jpi
#define __kJep1_
                  jpj
#define __kIe_
                  ntei
#define __kJe_
                 ntej
#define __kIep1_
                 nteip1
#define __kJep1_
                 ntejp1
                  __kIsm1_:__kIep1_,__kJsm1_:__kJep1_
#define IND 2D
```

These domain indices are declared as public variables in *par\_oce*. Their values are set by a new subroutine, *dom\_tile*, which takes a tile number as its argument. A negative or zero tile number is used to set the domain indices to the full domain. This is used to initialise the indices in *dom\_init*.

```
SUBROUTINE dom_tile(kntile)
  INTEGER , INTENT(in ) :: kntile
  INTEGER :: iitile, ijtile

IF( ln_tile .AND. kntile > 0 ) THEN
    iitile = 1 + MOD( kntile - 1, jpnitile )
    ijtile = 1 + (kntile - 1) / jpnitile

  ntile = kntile
  ntsi = 2 + (iitile - 1) * nn_tile_i
```

```
ntsj = 2 + (ijtile - 1) * nn_tile_j
      ntei = MIN(ntsi + nn_tile_i - 1, jpim1)
      ntej = MIN(ntsj + nn_tile_j - 1, jpjm1)
      ntsim1 = ntsi - 1
      ntsjm1 = ntsj - 1
      nteip1 = ntei + 1
      ntejp1 = ntej + 1
   ELSE
      ntile = 1
      ntsi = 2
      ntsj = 2
      ntei = jpim1
      ntej = jpjm1
      ntsim1 = 1
      ntsjm1 = 1
      nteip1 = jpi
      ntejp1 = jpj
   ENDIF
END SUBROUTINE dom tile
```

The tiling is implemented by looping over each tile at the time-stepping level (in *stp*) and calling *dom\_tile* to set the domain indices for the present tile subdomain. The indices are restored to the full domain after exiting the tiling loop.

```
DO jtile = 1, jpnijtile
   IF( ln_tile ) CALL dom_tile( jtile )
   CALL tra_ldf( kstp, Nbb, Nnn, ts, Nrhs )
END DO
IF( ln_tile ) CALL dom_tile( 0 )
```

tra\_ldf\_iso is then modified where necessary to work on the tile subdomain, e.g. local working arrays

```
REAL(wp), DIMENSION(jpi,jpj) :: zdkt, zdk1t, z2d
REAL(wp), DIMENSION(jpi,jpj,jpk) :: zdit, zdjt, zftu, zftv, ztfw
REAL(wp), DIMENSION(IND_2D) :: zdkt, zdk1t, z2d
REAL(wp), DIMENSION(IND_2D,jpk) :: zdit, zdjt, zftu, zftv, ztfw
```

and operations on global arrays.

Some code in *tra\_ldf* and *tra\_ldf\_iso*, mainly calls to other subroutines, has not been tiled and the best approach needs to be decided. This issue is summarised at the end of this document. At the moment this code is limited to running on one tile only; either the first (e.g. *WRITE* statements) or last (e.g. the *dia\_ptr\_hst* call) as appropriate.

### Alternative implementation using derived types

In this implementation, the domain indices are contained within a derived type *TILE*, declared in *dom\_oce*.

```
TYPE, PUBLIC :: TILE
```

```
INTEGER :: ntsi, ntsj, ntei, ntej, ntsim1, ntsjm1, nteip1, nteip1, ntile
END TYPE TILE
```

dom\_tile instead returns an instance of TILE

```
SUBROUTINE dom_tile(kntile, ktile)
  INTEGER , INTENT(in ) :: kntile
  TYPE(TILE), INTENT(out) :: ktile
  INTEGER :: iitile, ijtile

IF( ln_tile .AND. kntile > 0 ) THEN
    iitile = 1 + MOD( kntile - 1, jpnitile )
    ijtile = 1 + (kntile - 1) / jpnitile

  ktile % ntile = kntile
```

which is passed to *tra\_ldf\_iso* via an additional positional argument in *tra\_ldf* and *tra\_ldf\_iso* (shown below for the call to *tra\_ldf* in *stp*).

```
DO jtile = 1, jpnijtile
   CALL dom_tile( jtile, stile )
   CALL tra_ldf( stile, kstp, Nbb, Nnn, ts, Nrhs )
END DO
```

In do\_loop\_substitute.h90, duplicate DO loop macros have been added for tiled subroutines, e.g.

There is no need to call  $dom\_tile(0)$  to revert to the global domain or to initialise in  $dom\_init$ ; untiled subroutines simply do not have the additional positional argument. However, each call to  $tra\_ldf\_iso$  and its calling subroutines must have this additional argument, even if that subroutine has not yet been tiled. This results in a more widespread set of code changes.

Because *dom\_tile* returns an instance of *TILE*, rather than modifying public variables, the domain indices should be thread-private. However, this can also be achieved for the other approach by using an *OMP PRIVATE* directive.

Duplicate DO loop macros are needed in *do\_loop\_substitute.h90*, as not all subroutines will be tiled and *ktile* (the instance of *TILE* passed in) will be undefined in these cases.

In summary, using a derived type to pass around the domain indices is less tidy, involves more code changes and doesn't appear to have much functional benefit when compared to the approach using public variables.

## Things to consider

Untiled code within the tiling loop

A number of operations within the tiling loop will not function correctly, as they will now be called multiple times.

- 1. Operations that should only be performed once
  - WRITE statements, e.g. Iwp controlled writes to numout
  - NEMO timer (timing\_start)

- 2. Operations that should not proceed until all tiles are at the same point
  - MPI operations (*mpp\_sum*, etc)
  - Halo exchanges (*lbc\_lnk*)
  - NEMO timer (timing stop)
  - XIOS diagnostics (iom\_put)

In the trial implementation, operations under 1. and 2. are only performed for the first and last tiles respectively. This clearly will not work if the tiles are to be run concurrently using e.g. OpenMP.

Some formal solutions are being developed (XIOS diagnostics) and some may be quite straightforward (*lwp* can depend on the tile number, *ntile*, in addition to the processor number, *narea*). Other issues require more thought, particularly with regards to MPP.

### Extended haloes

The tiling implementation will be developed in tandem with the extended haloes development. There are a few potential issues here.

- 1. Refactoring of loops
- 2. Movement of *lbc\_lnk* calls
- 3. Use of extended haloes, nn\_hls > 1

The trial implementation of tiling mainly involves higher level code changes via the DO loop macros and only changes lower level code in a few cases, so 1. will hopefully not be too disruptive.

The tiling development should be able to proceed on 1 CPU, so 2. should not be an issue for the time being. Do we expect any *lbc\_lnk* calls to remain in the tiling loop after the extended halo development? If so, how should we deal with this when using MPP?

Potentially, resolving 3. could be as simple as modifying the calculations where  $nn_hls = 1$  is assumed, e.g. the domain indices in  $dom\ tile$ 

```
ntsi = 1 + nn_hls + (iitile - 1) * nn_tile_i
ntei = MIN(ntsi + nn_tile_i - 1, jpi - nn_hls)
ntsi_1 = ntsi - nn_hls
ntei 1 = ntei + nn hls
```

where *ntsi\_1* and *ntei\_1* have replaced *ntsim1* and *nteip1* as the indices for the full domain in this example.